## The Envelope Method

## for computing orthogonal eigenvectors belonging to isolated clusters of very close eigenvalues of a symmetric tridiagonal matrix

B. N. Parlett<br>Math. Dept. and Comp. Sci. Div. of the EECS Dept. University of California, Berkeley<br>parlett@math.berkeley.edu

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## Eigenvectors - The pure mathematician's solution

Givens' method

$$
(T-\lambda /) \mathbf{x}=\mathbf{0}
$$

$$
\left[\begin{array}{ccccc}
\alpha_{1}-\lambda & \beta_{1} & & & \\
\beta_{1} & \alpha_{2}-\lambda & \beta_{2} & & \\
& \beta_{2} & \alpha_{3}-\lambda & \beta_{3} & \\
& \ddots & \ddots & \ddots & \\
& & \beta_{n-2} & \begin{array}{c}
\alpha_{n-1}-\lambda \\
\beta_{n-1}
\end{array} & \beta_{n-1} \\
& & & \alpha_{n}-\lambda
\end{array}\right]\left[\begin{array}{c}
x_{1} \\
x_{2} \\
x_{3} \\
\vdots \\
x_{n-1} \\
x_{n}
\end{array}\right]=\mathbf{0}
$$

Assumption: $\beta_{i} \neq 0$, all $i$
[for the all talk]

Theorem 1. $\beta_{i} \neq 0$, all $i \Longrightarrow$ eigenvalues distinct
Theorem 2. $\beta_{i} \neq 0$, all $i \Longrightarrow x_{1} x_{n} \neq 0$
Givens method: set $x_{1}=1$

| Eqn. 1 | determines |
| :---: | :---: |
| Eqn. 2 | $x_{2}$ |
|  |  |
| $x_{3}$ |  |

$$
\begin{aligned}
& \text { Eqn. } n-1 \quad " \quad x_{n} \\
& \text { [Eqn. } n \text { is redundant] } \\
& \text { Normalize, if required } \\
& \left\langle\text { could start with } x_{n}=1 \text { and proceed in reverse }\right\rangle
\end{aligned}
$$

Properties. Perfect in exact arithmetic with exact $\lambda$.
Defect 1 . What if 2 eigs agree to working precision?
What if 4 eigs agree too working precision?
Defect 2. Can fail even for an isolated eigenvalue ( $W_{21}^{-}$)

# Is it difficult to compute numerically orthogonal eigenvectors belonging to an isolated cluster of very close eigenvalues? 

Ans: yes, if you compute them one by one.

Ans: no , if you compute them all together.

## Envelopes

The envelope of a vector $v$ is

$$
|v|=\left[\left|v_{1}\right|,\left|v_{2}\right|, \ldots,\left|v_{n}\right|\right]^{t}
$$

Let $Q$ be any orthonormal basis for a subspace $S \subseteq \mathbb{R}^{n}$ of dimension $p$. Then its envelope $\mathcal{E}_{s}$ is given by

$$
\begin{aligned}
& \mathcal{E}_{s}(i)=\|Q(i, 1: p)\|_{2}, \quad i=1,2, \ldots, n \\
& \left\|\mathcal{E}_{s}\right\|_{2}=\|Q\|_{F}=\sqrt{p}
\end{aligned}
$$

The envelopes of invariant subspaces belonging to isolated clusters of very close eigenvalues of tridiagonal $T$ have hills and valleys.

The smaller is

## cluster width

cluster gap
the deeper are the valleys.
The number of hills $\leq$ the number of eigenvalues in the cluster.

## Distinguished sparse basis

To each hill associate a vector that is zero except for the hill (suitably extended).
Some hills may have two vectors assigned to them.


Figure: An eigenvector of a tridiagonal: most of its entries are negligible

Matrix comes from nuclear chemistry
George Fann matrix, $n=966$


Figure: Snapshot of Envelope (108 eigenvalues)
$n=2053$, cluster size 108
Cluster determined by submatrix 1:515

## How to construct a distinguished basis from the envelope

Take each hill down to the valley on each side and extend smoothly to zero to obtain the indices first and last


This gives one submatrix $T$ (first: last) per hill.

In pratice, we will not have the envelope.
We want to create the index pair (first:last) from the tridiagonal itself.

There are several methods.
We will describe a new inexpensive one later.

To create the basis vector(s)

- Use the submatrix $T$ (first.last)
- Compute the eigenpair(s) with eigenvalue in the cluster interval.

| $\lambda_{21}$ | $10.746194182903393 .$. |
| :---: | :---: |
| $\lambda_{20}$ | $10.746194182903322 .$. |
| $\cdot$ | $\cdot$ |
| $\cdot$ | $\cdot$ |
| $\lambda_{13}$ | 6.0002340. |
| $\lambda_{12}$ | $6.0002175 .$. |
| $\cdot$ | $\cdot$ |
| . | $\cdot$ |
| $\lambda_{3}$ | $0.9475 .$. |
| $\lambda_{2}$ | $0.2538 .$. |
| $\lambda_{1}$ | $-1.1254 .$. |

Table: Selected eigenvalues of $W_{21}^{+}$

$$
W_{21}^{+}=\left[\begin{array}{ccccccccccc}
10 & 1 & & & & & & & & & \\
1 & 9 & 1 & & & & & & & & \\
& 1 & 8 & 1 & & & & & & & \\
& & \ddots & \ddots & \ddots & & & & & & \\
& & & 1 & 1 & 1 & & & & & \\
& & & & 1 & 0 & 1 & & & & \\
& & & & & 1 & 1 & 1 & & & \\
& & & & & & \ddots & \ddots & \ddots & & \\
& & & & & & & 1 & 8 & 1 & \\
& & & & & & & & 1 & 9 & 1 \\
& & & & & 10
\end{array}\right]
$$



Figure: Envelope for $\lambda_{12}, \lambda_{13}$ from $W_{21}^{+}$


Figure: Vectors $z_{+}$and $z_{-}$for the pair near 6 on a log scale
$0<x<1, \quad x \longrightarrow-1 / \log _{10} x, \quad$ correct sign normalized


Figure: Bisectors of $z_{+}$and $z_{-}$on a log scale




Figure: Basis eigenvectors of $W_{100}$ for $\left\{\lambda_{93}, \lambda_{94}, \lambda_{95}\right\}$
Glued Wilkinson matrix

$$
G_{2}\left(W_{25}^{+}, 4,0.3\right)
$$

4 copies of $W_{25}^{+}$, glue $=0.3$
Cluster structure 3-2-3
12.5778... 3
12.74619 .

2
$12.93911 \ldots 3$



Figure: Basis eigenvectors of $W_{100}$ for $\left\{\lambda_{96}, \lambda_{97}\right\}$




Figure: Basis eigenvectors of $W_{100}$ for $\left\{\lambda_{98}, \lambda_{99}, \lambda_{100}\right\}$


Figure: $G_{2}\left(W_{201}^{+}, 2, \sqrt{\varepsilon}\right), \quad \lambda_{3}, \lambda_{4} \sim 0.25$


Figure: Eigenvectors $x_{79}, x_{80}, x_{81}, x_{82}$ for $T_{87}$ on a log scale

Wanted:
a cheap way to approximate the submatrices

## Results complementary to Gersgorin's Circle Theorem

$B \in \mathbb{C}^{n \times n}$

$$
G_{j}^{\text {row }}:=\left\{\zeta:\left|\zeta-b_{j j}\right| \leq \sum_{k \neq j}\left|b_{j k}\right|\right\}
$$

G's Theorem. Each $\lambda \in \operatorname{eig}(B)$ is located in at least one $G$-disk.

## A complementary result

If

$$
\lambda \notin G_{j}^{\text {row }} \text { for } j=p: q, \quad p<q,
$$

then
$\lambda$ 's column eigenvector decays in entries $p: q$.
(the direction of decay varies, sometimes from $p$ to $q$, sometimes $q$ to $p$ )

Corollary. The (local) maximal elements of an eigenvector occur only for indices $k$ for which $\lambda \in G_{k}$.

For large $n$ this corollary makes searching for a maximal entry more efficient.
E.g.

If $\lambda \notin G_{i}, i=1, \ldots, 500$, then start computing at index 501.

Are there any uses for indices $i$ for which $\lambda \notin G_{i}$ ?

## Symmetric Tridiagonal Case

$$
T=\left[\begin{array}{cccccc}
\alpha_{1} & \beta_{1} & & & & \\
\beta_{1} & \alpha_{2} & \beta_{2} & & & \\
& \beta_{2} & \alpha_{3} & \beta_{3} & & \\
& & \ddots & \ddots & \ddots & \\
& & & \beta_{n-2} & \alpha_{n-1} & \beta_{n-1} \\
& & & & \beta_{n-1} & \alpha_{n}
\end{array}\right]
$$

Suppose that

$$
T-\lambda I=L D L^{t}
$$

$L:=\left[\begin{array}{ccccc}1 & & & & \\ l_{1} & 1 & & & \\ & l_{2} & 1 & & \\ & & \ddots & \ddots & \\ & & & I_{n-1} & 1\end{array}\right], \quad D:=\left[\begin{array}{lllll}d_{1} & & & & \\ & d_{2} & & & \\ & & \ddots & & \\ & & & d_{n-1} & \\ & & & & 0\end{array}\right]$
$(T-\lambda I) z=\mathbf{0} \Longrightarrow L^{t} z=\mathbf{e}_{n} z_{n} \Longrightarrow z_{j}=-l_{j} z_{j+1}, j=n-1, \ldots, 1$

Lemma. If $\lambda \notin G_{j}$ and $\left|I_{j-1}\right|<1$ then $\left|\ell_{j}\right|<1$.
Proof. $\quad \iota_{j}=\frac{\beta_{j}}{d_{j}}$

$$
d_{j}=\alpha_{j}-\lambda-l_{j-1} \beta_{j-1}
$$

$$
\begin{aligned}
\left|d_{j}\right| & \geq\left|\alpha_{j}-\lambda\right|-\left|I_{j-1}\right|\left|\beta_{j-1}\right| \\
& >\left|\alpha_{j}-\lambda\right|-\left|\beta_{j-1}\right|, \quad \text { since }\left|I_{j-1}\right|<1
\end{aligned}
$$

$$
\left|I_{j}\right|<\frac{\left|\beta_{j}\right|}{\left|\alpha_{j}-\lambda\right|-\left|\beta_{j-1}\right|}
$$

$$
<1, \quad \operatorname{since} \lambda \notin G_{j} \square
$$

Theorem. If $\lambda \notin G_{j}$ for $j=p: q, p<q$, and $\left|I_{p-1}\right|<1$ then

$$
\left|f_{j}\right|<1 \quad \text { for } j=p-1: q
$$

and

$$
\left|z_{j}\right|<\left(\prod_{i=j}^{q}\left|\ell_{i}\right|\right)\left|z_{q+1}\right| \quad \text { for } j=p-1: q
$$

Similar results for $T-\lambda I=U D U^{t}$.
Heuristic: replace $d_{i}^{+}$and $d_{i}^{-}$by $\alpha_{i}-\lambda$ Hence

$$
\prod_{i=j}^{q}\left|l_{i}\right| \approx \frac{\prod_{i=j}^{q} \beta_{i}}{\prod_{i=j}^{q}\left(\alpha_{i}-\lambda\right)}=: \frac{\text { num }}{\operatorname{den}}
$$

Test: num $\leq$ tol.den

$$
\text { tol }=10^{-12} \text { or } 10^{-15}
$$

## General Case

$$
\begin{gathered}
B-\lambda I=L^{+} D^{+} U^{+} \\
u_{k j}^{+}=\frac{b_{k j}}{d_{k}^{+}}, \quad k=1,2, \ldots, j-1 \\
d_{j}^{+}=\left(b_{j j}-\lambda\right)-\sum_{k<j} b_{j k} u_{k j}^{+} \\
u_{j l}^{+}=\frac{b_{j l}}{d_{j}^{+}}, \quad I=j+1: n \\
\sum_{l>j}\left|u_{j l}^{+}\right|=\frac{1}{\left|d_{j}^{+}\right|} \sum_{l>j}\left|b_{j l}\right|
\end{gathered} \begin{gathered}
\leq \frac{\sum_{l>j}\left|b_{j l}\right|}{\left|b_{j j}-\lambda\right|-\left(\sum_{k<j}\left|b_{j k}\right|\right) \max _{k<j}\left|u_{k j}^{+}\right|} \\
<\frac{\sum_{l>j}\left|b_{j l}\right|}{\left|b_{j j}-\lambda\right|-\sum_{k<j}\left|b_{j k}\right|}, \text { if } \max _{k<j}\left|u_{k j}^{+}\right|<1 \\
<1, \text { if } \lambda \notin G_{j}^{r o w}
\end{gathered}
$$

Theorem. If $\max _{k<p}\left|u_{k p}^{+}\right|<1$ and $\lambda \notin G_{j}^{\text {row }}$ for $j=p: q$ then

$$
\sum_{l>m}\left|u_{m l}^{+}\right|<1 \text { for } m=p: q .
$$



Figure: $W_{51}^{+}, \quad \lambda_{21}$

True vector vs. heuristic, $n=51, \lambda=15$


Figure: $W_{51}^{+}, \quad \lambda_{31}$


Figure: $W_{51}^{+}, \quad \lambda_{41}$

True vector vs. heuristic, $n=51, \lambda=25.7462$


Figure: $W_{51}^{+}, \quad \lambda_{51}$

## Conclusion

For symmetric tridiagonal matrices the computation of numerically orthogonal eigenvectors for

> isolated cluster of close eigenvalues
is
easy and rapid
if the cluster is treated as a whole and

> almost impossible
if they are computed one by one.

